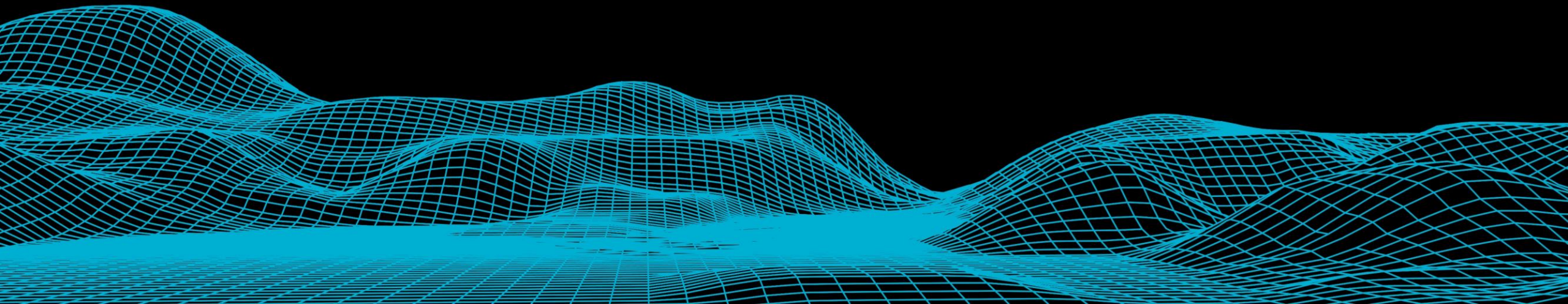




# Performance Optimization of the CMAQ model on Microsoft Azure

Steve Roach

Azure / HPC Technical Specialist  
[sroach@microsoft.com](mailto:sroach@microsoft.com)



# Agenda

1. Describe Azure Platform
2. Discuss Factors which impact application performance
  - a) Process Pinning
  - b) Storage

# Azure: cloud built for HPC & AI

Genuine HPC & AI approach:  
platforms, benchmarks, people

Purpose-built hardware for the best  
performance, optimized price-  
performance and differentiated  
solutions

TTM availability of leading hardware  
innovations to accelerate “time to  
solutions” for customer workloads

Strategy to leverage leading internal  
production workloads using the same  
systems for mission critical offerings on  
Azure.



Supercomputing  
scale for the  
most demanding  
applications

InfiniBand  
HPC & AI  
clusters for best  
performance on  
real workloads

Compute  
optimized VMs  
with low latency  
ethernet

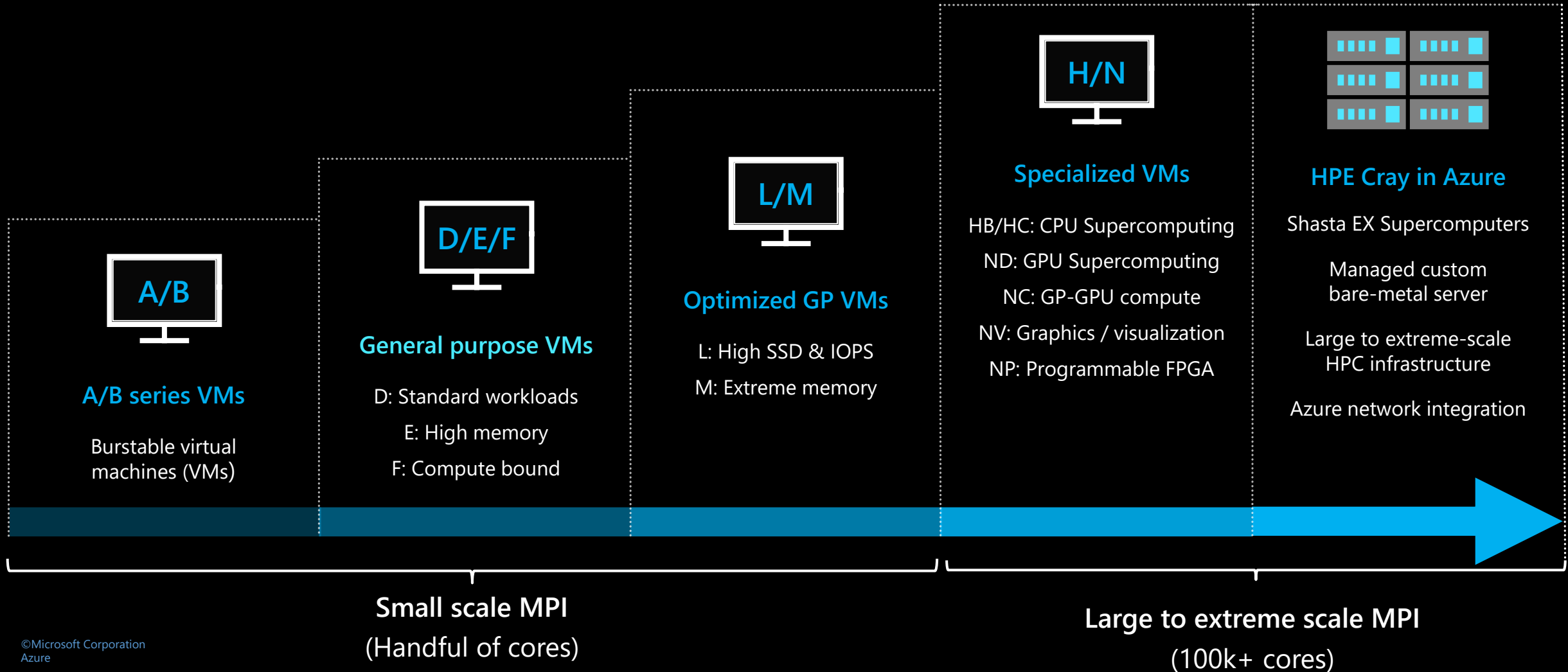
Azure

Azure is  
the only  
public cloud  
provider  
offering the  
full range  
of HPC and AI  
capabilities

Compute  
optimized VMs  
with low latency  
ethernet

Other clouds

# Azure looks a lot like a HPC Datacenter



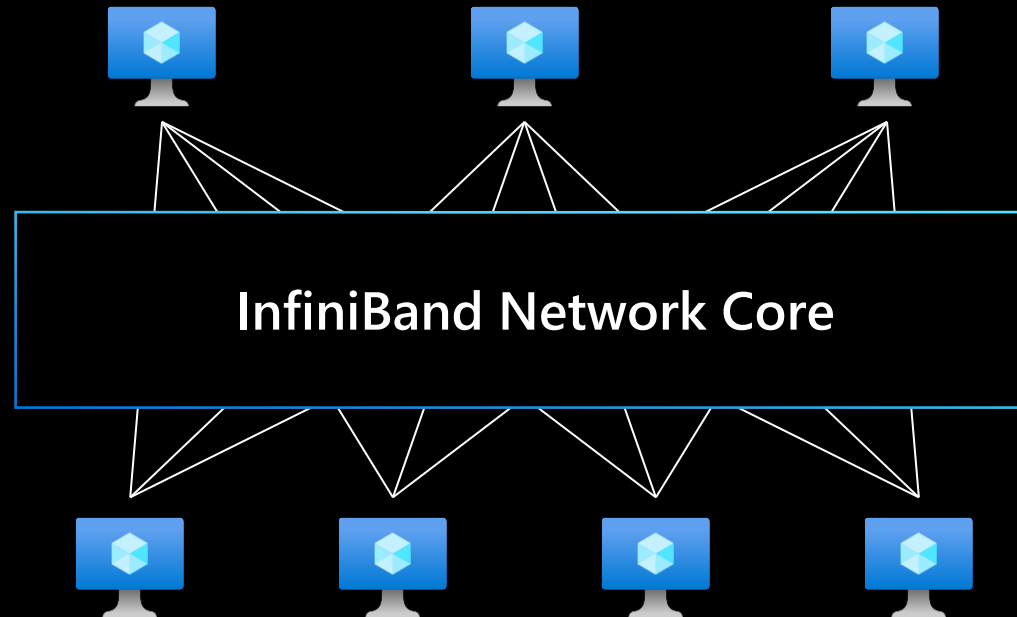
Non-blocking Fat Tree topology

Hardware offload of MPI collectives

Full MPI & NCCL Integration

< 1.5 microsecond latencies

Up to 1.6 Tb/s per VM



Bare-metal passthrough

Intelligent Adaptive Routing

Dynamic Connected Transport

# SR-IOV Goodness on Azure

## MPICH Derivatives:

- ✓ MPICH
- ✓ IntelMPI
- ✓ MVAPICH2
- ✓ Microsoft MPI
- ...

## OpenMPI Derivatives:

- ✓ OpenMPI
- ✓ HPC-X
- ✓ Platform MPI
- ...

- Feature parity with bare metal
- Prior to SR-IOV enablement Supported only IntelMPI (dapl) and MSMPI

# HBv3 Upgrade

Milan-X comes to Azure HPC

HBv3 Virtual Machines enhanced with AMD 3<sup>rd</sup> Gen EPYC with 3D v-cache

No customer/partner changes required, same HBv3 VM sizes

3x L3 cache increase per core, chiplet, socket, and VM (1.5 GB)

Accelerates HPC applications bound by memory performance

Increases effective memory bandwidth up to ~630 GB/s

Decreases effective memory latency by as much as 51%



Performance & Scalability of HBv3 VMs with Milan-X CPUs



# What is Milan-X and how does it affect performance?

Architecturally, Milan-X differs from Milan only by virtue of having 3x as much L3 cache memory per core, CCD, socket, and server.

CPU	Xeon 2690 v4 "Broadwell"	Xeon Gold 6148 "Skylake"	Xeon 8280 "Cascade Lake"	EPYC 7742 "Rome"	EPYC 7V73X "Milan-X"
Cores/2S server	28	40	56	128	128
L3 cache/2S server	70 MB	55 MB	77 MB	512 MB	1,536
Relative size	1x	0.8x	1.1x	7.3x	22x

Examples of workloads that can benefit from larger L3 cache are:

- Computational fluid dynamics (CFD) – memory bandwidth
- Weather simulation – memory bandwidth
- Explicit finite element analysis (FEA) – memory bandwidth
- EDA RTL simulation – memory latency



# Advantages of Fewer Cores per Node

	120 Cores	96 Cores
RAM per Core	3.75 GB	4.67 GB
Memory B/W per Core	2.91 GB/s	3.65 GB/s
L3 Cache per Core	12.8 MB	16 MB

# Importance of Process Pinning

For MPI applications, optimal pinning of processes can lead to significant application performance improvements for under subscribed systems.

In the chiplet design, AMD has essentially integrated a bunch of smaller CPUs together to provide a socket with 64 cores (8 - 16 smaller CPUs with 4-8 cores each).

**To maximize the performance from each core it is important to balance the amount of L3 cache and memory bandwidth per core.**

# Pinning Example

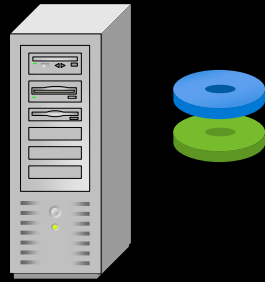
```
setenv PIN_PROCESSOR_LIST "--bind-to cpulist:ordered --cpu-set  
0,1,2,3,4,5,8,9,10,11,12,13,16,17,18,19,20,21,24,25,26,27,28,29,30,31,32,33,34,35,38,3  
9,40,41,42,43,46,47,48,49,50,51,54,55,56,57,58,59,60,61,62,63,64,65,68,69,70,71,7  
2,75,76,77,78,79,80,81,84,85,86,87,88,89,90,91,92,93,94,95,98,99,100,101,102,103,1  
06,107,108,109,110,111,114,115,116,117,118,119 --report-bindings "
```

```
( /usr/bin/time -p mpirun -np $NPROCS $PIN_PROCESSOR_LIST --rank-by slot -  
mca coll ^hcoll -x LD_LIBRARY_PATH -x PATH -x PWD $BLD/$EXEC ) |& tee  
buff_${EXECUTION_ID}.txt
```

# Storage Options



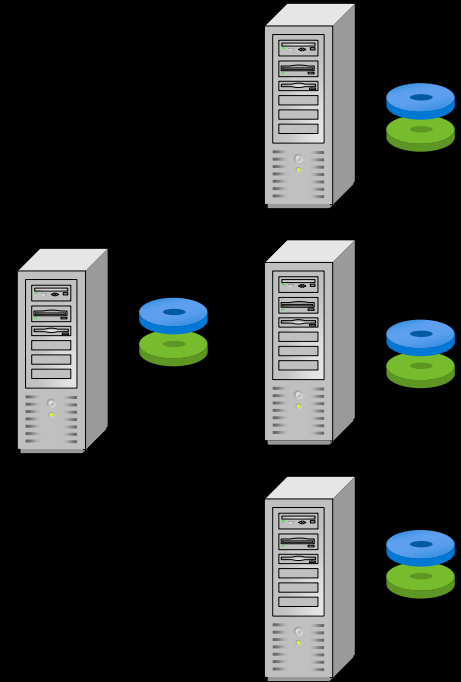
Azure Files  
with NFS



Stand Alone  
NFS Server



Azure NetApp  
Files



Lustre Cluster

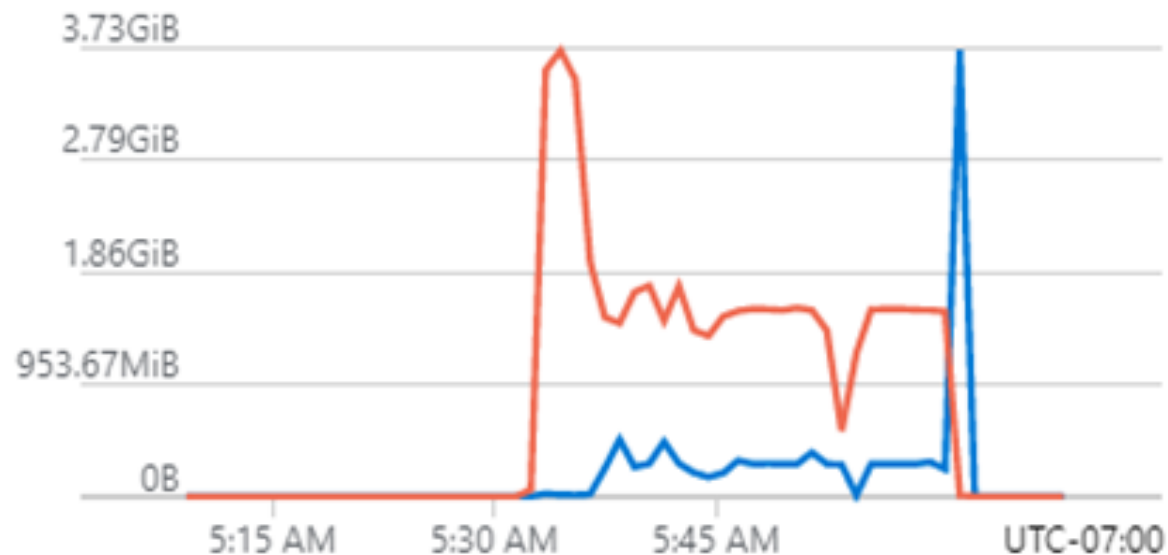
# Azure HPC File System Portfolio

	NFS Blobs*	NFS Files*	HPC Cache	Azure NetApp Files	*aaS PVFS	Cray ClusterStor
File System	NFSv3	NFSv4.1	NFSv3 NFSv4.1* SMB2.1*	NFSv3 NFSv4.1 SMB	Lustre BeeGFS	Lustre
Characteristics	Large size Medium throughput Sequential access Read-heavy	Medium size Medium throughput Random access In-place updates	Large file Sequential reads Optimizes latency and throughput Caches multiple source NAS environments	Medium size Medium throughput Random access In-place updates Low latency	Large size High throughput Sequential access 10-15 MiB/s per provisioned TiB	Very large size High throughput Sequential access 10-15 MiB/s per provisioned TiB
Use cases	Legacy NFS apps Backup and archive Analytics	Shared app data Databases Container storage Home directories	HPC up to 8GB/s Read heavy Cloud burst from on-prem NAS Multi-source file system (on-prem and in-cloud)	Enterprise application migrations NFS home dirs	Built to specs Durable or Ephemeral options	Long-lived HPC jobs (weeks/months) Data protection and tiering in Blob storage
File system details	5 PiB 100K IOPS (premium) 12.5GB/s throughput	100TiB 100K IOPS 10GB/s	50TiB cache 240k IOPS 8 GB/s	100 TiB 300K IOPS 4.5GB/s	<1PB TBD IOPS up to 200GB/s	5PiB, 15PiB, 45 PiB 240k IOPS 200GB/s
Region availability	Broad	Broad	Broad	Select	Broad	On-demand

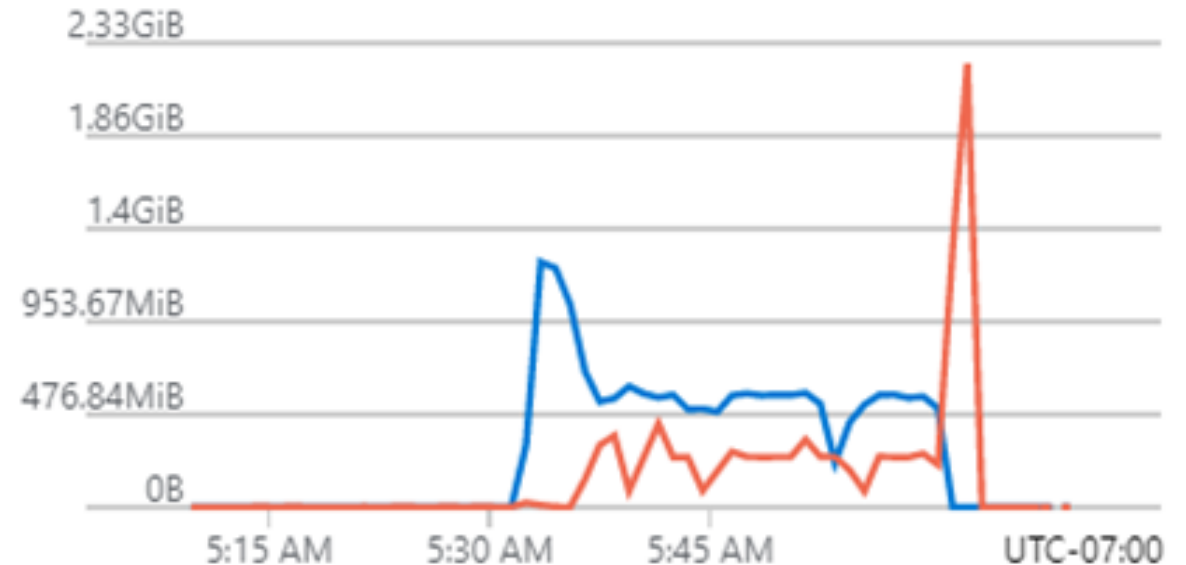
# CMAQ I/O Activity

## Stand Alone NFS Server

Network (total)

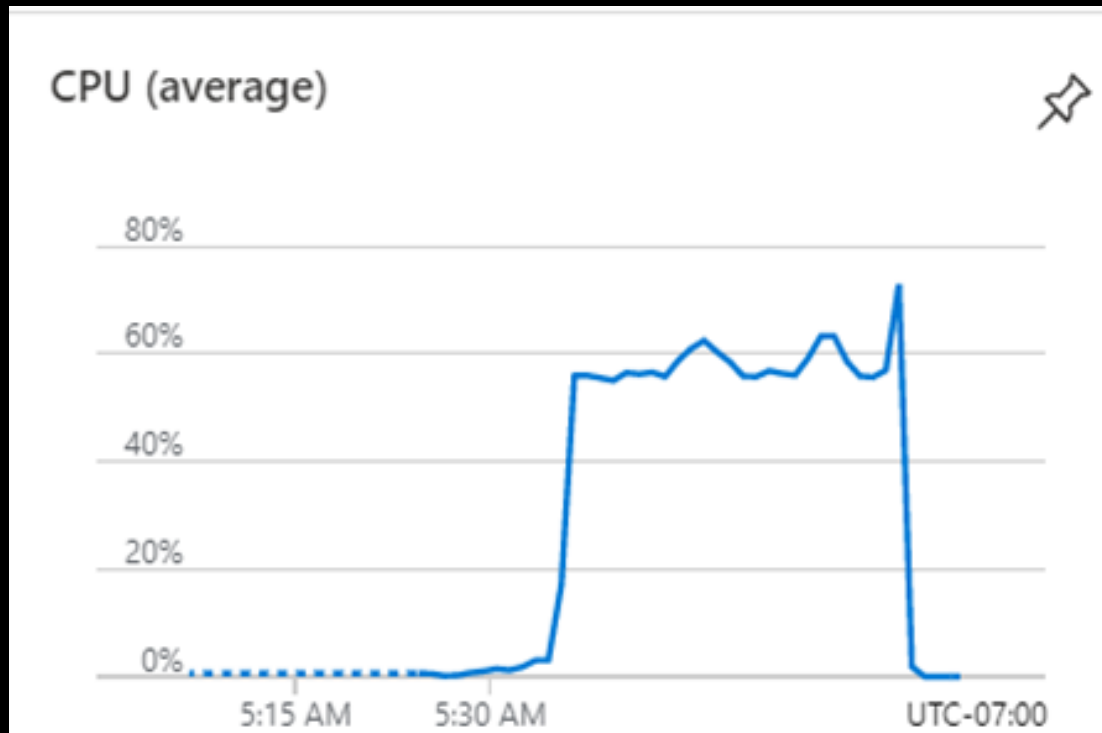


Disk bytes (total)

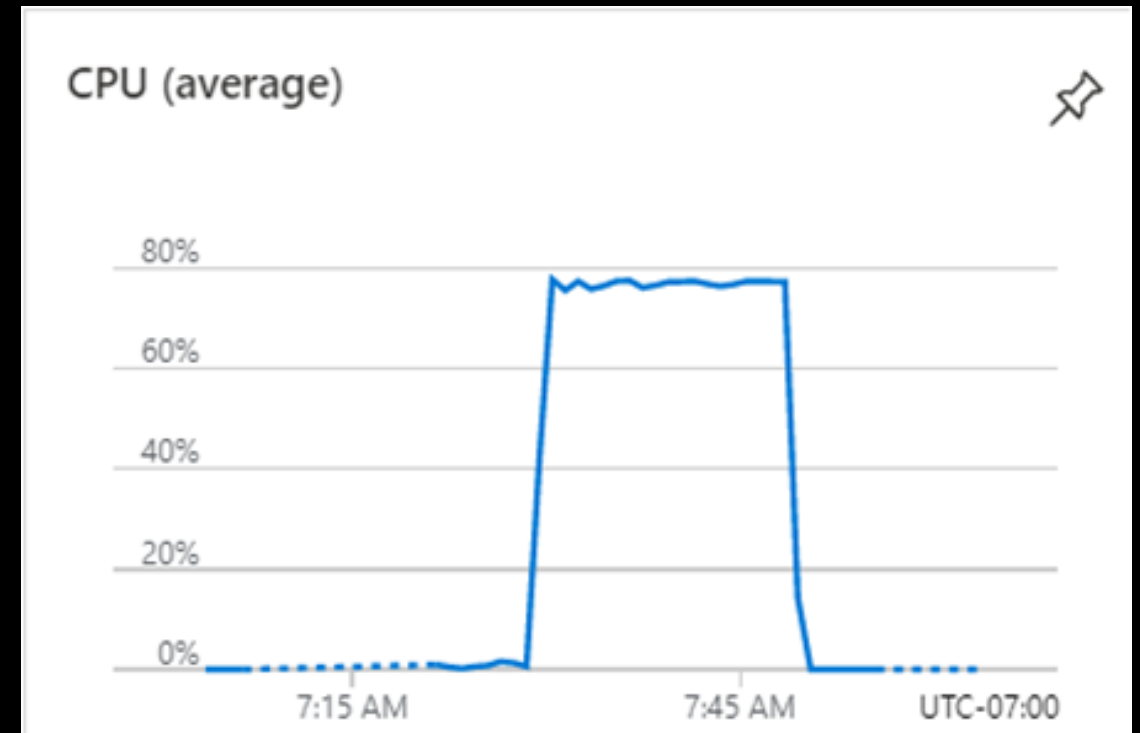


# CMAQ CPU Activity – Average across 3 Compute Nodes

Stand Alone NFS Server

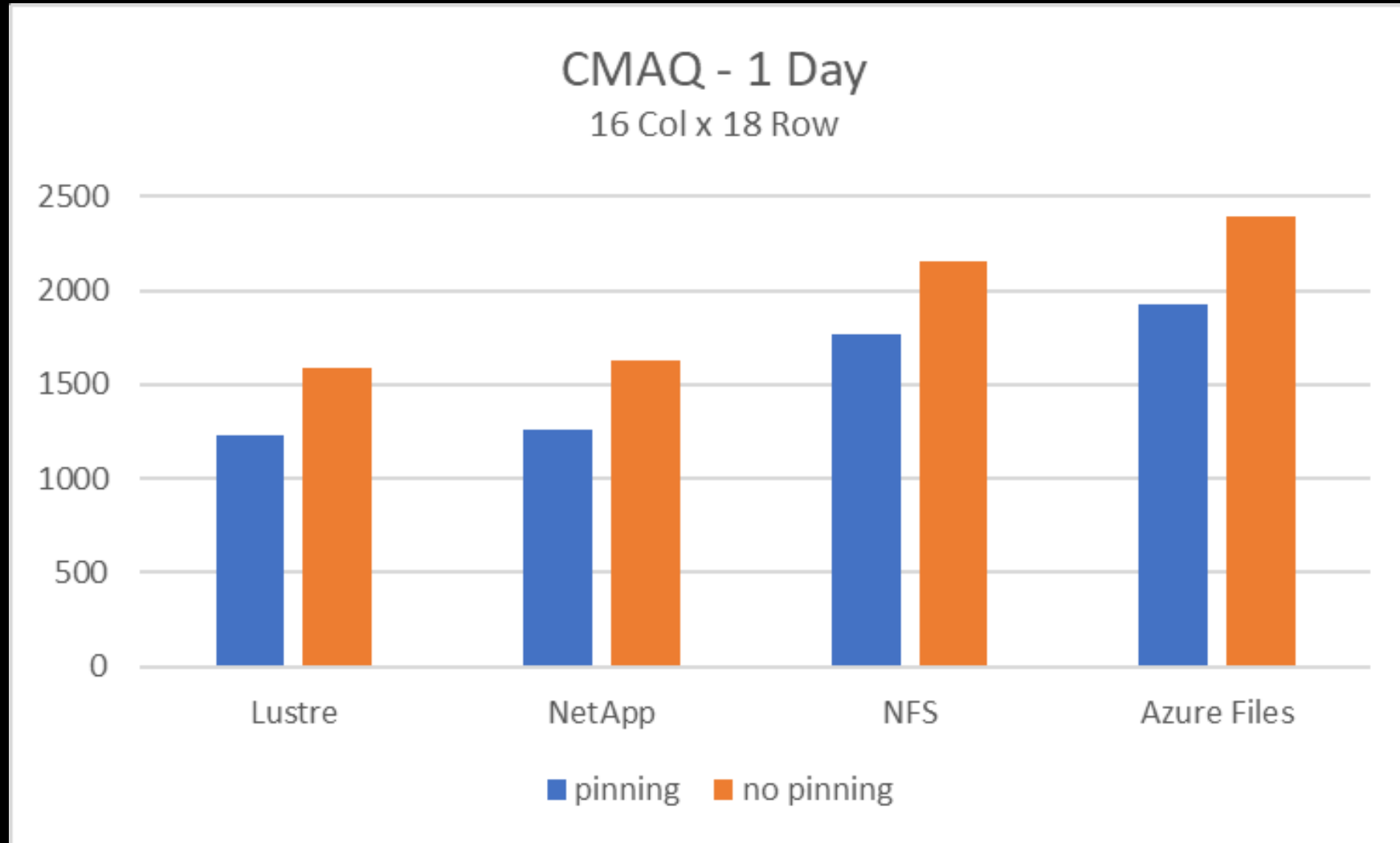


Lustre





# Test Results – 3 nodes, 96 ppn



A photograph of Earth from space, showing the curvature of the planet and city lights at night. The word "Questions?" is overlaid in white text.

# Questions?